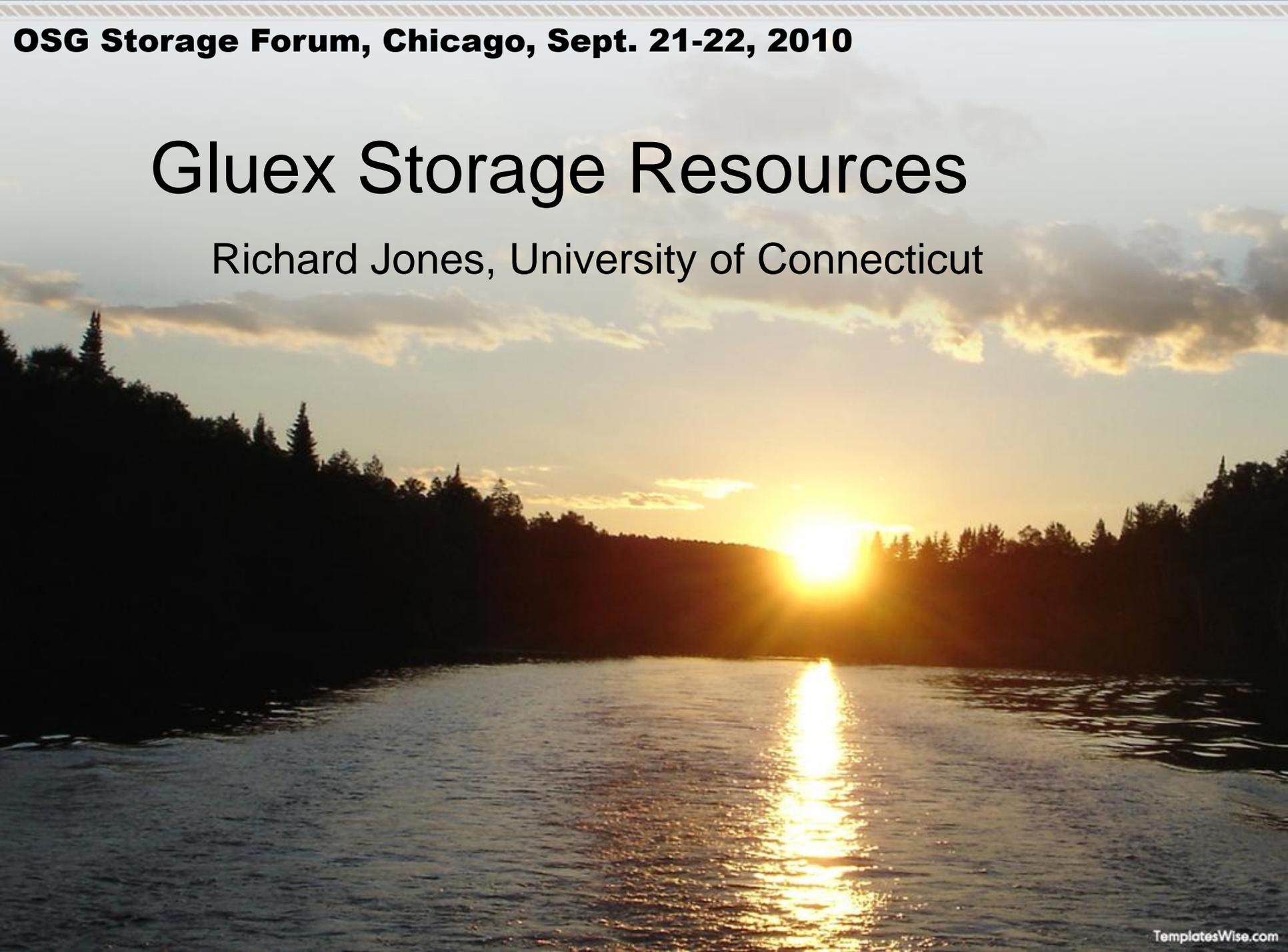


# Glutex Storage Resources

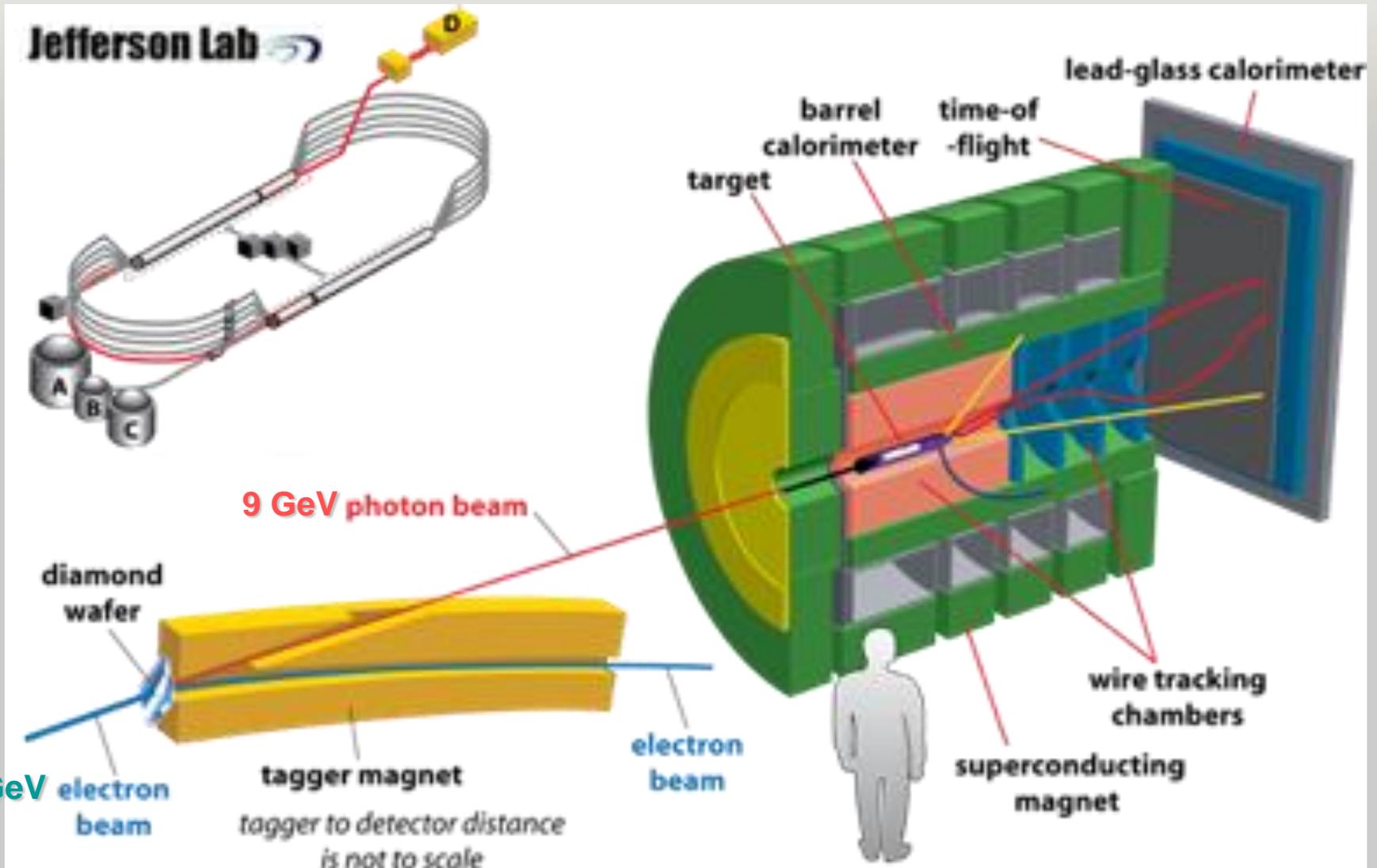
Richard Jones, University of Connecticut



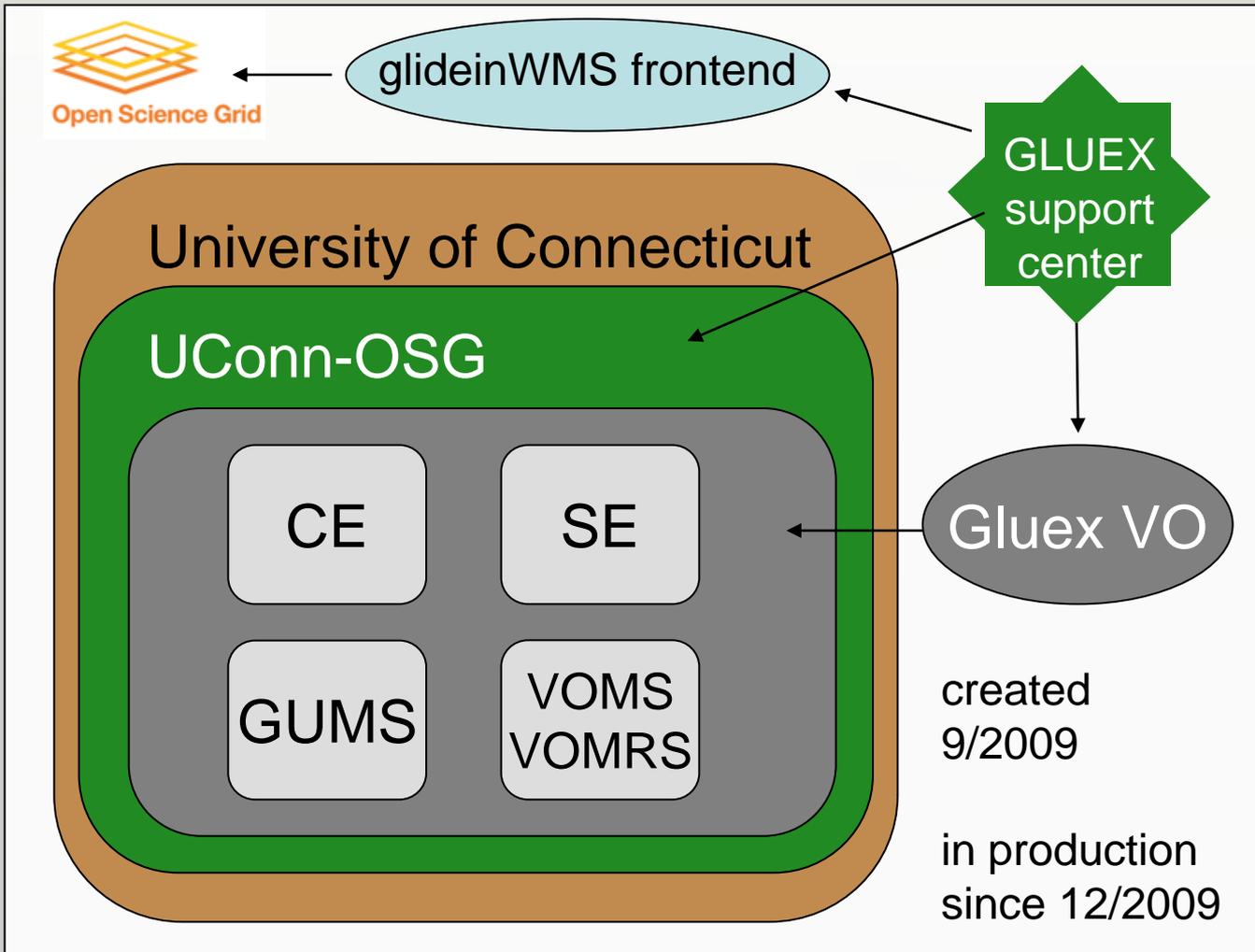
# Outline

- who is Gluex?
- data storage and delivery needs
- existing resources, experience
- plans and outlook

# GlueX – search for hybrid mesons



# GlueX VO – the collaboration



## member institutions

- University of Athens
- Carnegie Mellon Univ.
- Catholic University
- Christopher Newport Univ.
- University of Connecticut
- Florida International Univ.
- Florida State University
- University of Glasgow
- IHEP Protvino
- Indiana University
- Jefferson Lab
- U. of Massachusetts
- North Carolina A&T State
- U. of North Carolina
- Santa Maria University
- University of Regina

# GlueX VO – data storage, delivery needs

- raw data: 10 kB/event, 20 kHz event rate = 2 TB / year
  - ◆ archived on Jlab site (tape library)
  - ◆ reconstruction -> DST with 5% of raw events, 20 kB/event

**200 TB/year, 5 years = 1 PB total for export offsite**

- Monte Carlo: 20 kB/event, 100 kB/s on a 2.5GHz core2
  - ◆ minimum-biased event sample most challenging
  - ◆ ideally should approach raw data statistics
  - ◆ simulate, reconstruct, keep only MC DST

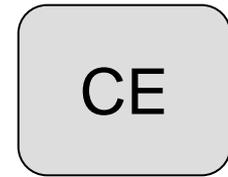
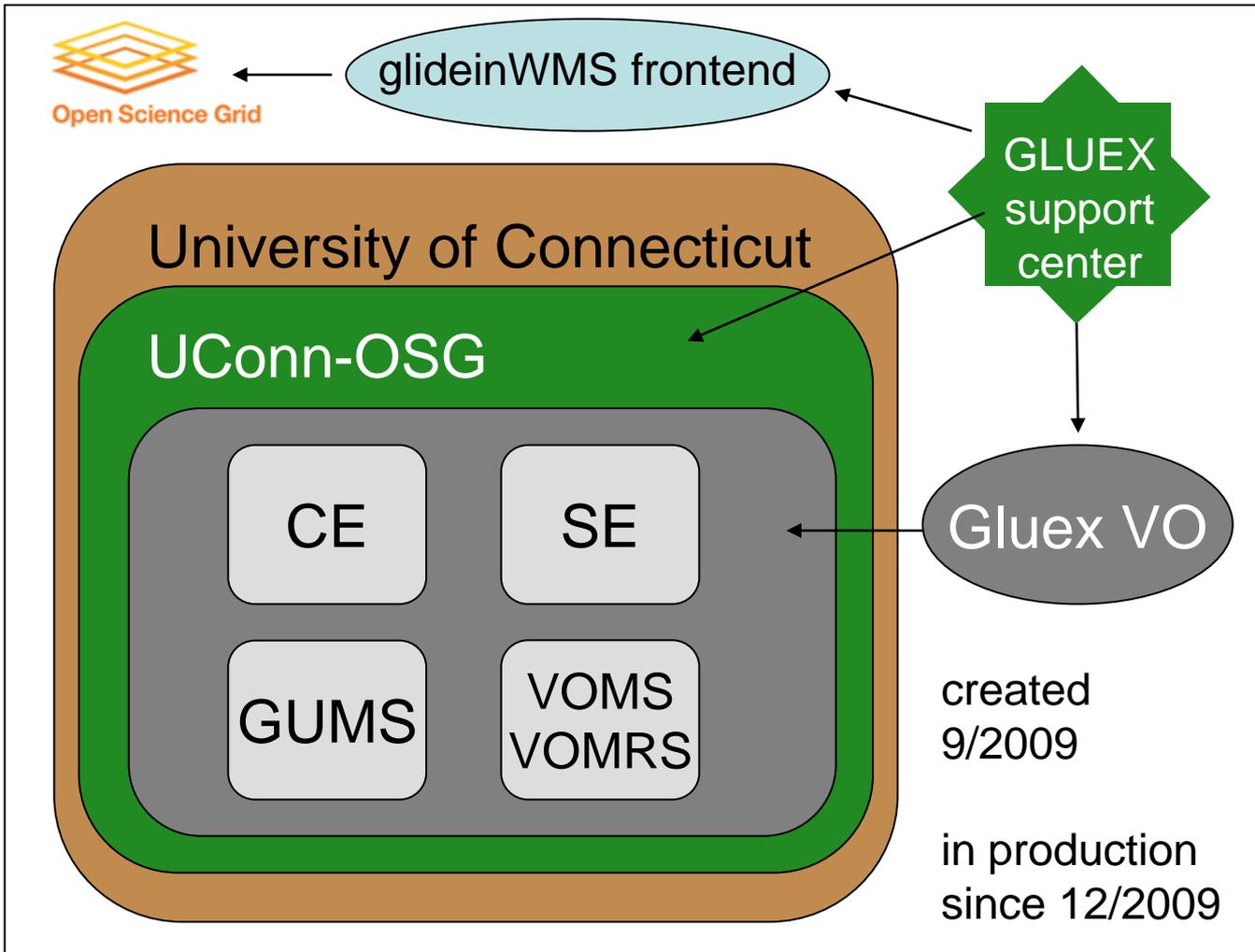
**100 TB/year, 5 years = 500 TB for provision offsite**

- ◆ production targeted for OSG (min.bias sample: *30M hours*)

# GlueX VO – data storage, delivery needs

- analysis: cuts to select exclusive final states
  - ◆ reduction jobs go to where data resides (major sites)
  - ◆ micro-DSTs (root trees, few TB each) per analysis
  - ◆ Monte Carlo (not min.bias) needed on-demand
  
- PWA fits: performed on dedicated GPU hardware
  - ◆ should be interactive
  - ◆ requires real-time access to micro-DSTs
  - ◆ may move toward scheduled GPU resources

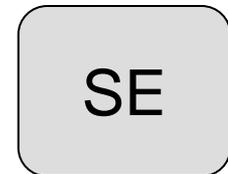
# GlueX VO: existing resources



condor-jobmanager

280 x86\_64 cores

+ 100 xeon cores



dcache (w/o HSM)

30 pool nodes

70 pools

# GlueX VO: why dcache?

## ❑ **experience before dcache (pre 2004)**

- **pvfs** – parallel virtual filesystem (*R. Ross, Clemson*)
- 10K files (2-3 TB) splintered across 15 nodes
- performance ok (could saturate network)
- administration painful: **1 server down => file system hangs**
- **kernel integration: encumbers OS upgrade scheduling**
- **metadata uncopyable, files unrecoverable if corrupted**
- **zero redundancy, frequent data loss**

## ❑ **dcache seemed to answer many of these problems**

- layered on top of an ordinary unix filesystem
- uses the built-in kernel nfs support (no custom kernel modules)
- metadata stored in a standard database
- filesystem robust against single pool node failures

# GlueX VO: why dcache?

## ❑ **experience with dcache (2004-2009, pre-OSG)**

- peak performance somewhat worse than pvfs (factor 2-3)
- net throughput with parallel jobs was about equal to pvfs
- overall experience was much, much better
  - rare data loss (3-4 times in 5 years, human error)
  - robust hands-off operation for weeks at a time (~1TB i/o per wk)
  - stable across OS upgrades

## ❑ **recent experience (with operation as a OSG SE)**

- requires considerable work to keep it running
- suffers from an authentication bottleneck (GUMS timeouts)
- seeing out-of-heap-memory errors under heavy load
- SRM response seems sluggish (30s for a short ls)
- first time full authentication layers are exercised

# GlueX Storage: Plans and Outlook

## Why dcache might work for us:

1. the right mix of protocols: SRM/gridftp, xrootd, plain http-get
2. flexible configuration with control over replica management
3. nfs namespace introspection
4. ongoing development, large user base
5. no kernel-space code

## Why dcache might not work for us:

1. authentication/authorization performance
2. SRM transaction overhead
3. lack of a comprehensive “fsck” tool
4. pain of administration (robustness under realistic conditions)

***Next on our list to evaluate: hadoop***